

# プログラミング言語(SICP) 3. Modularity, Objects, and State 3.5 streams

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## 5月16日・本日のメニュー

6月13日(木) 中間テスト

範囲は第3章

3.5 Streams

3-5-1 Streams are Delayed Lists

3-5-2 Infinite Streams

3-5-3 Exploiting the Stream Paradigm

3-5-4 Streams and Delayed Evaluation

3-5-5 Modularity of Functional Programs  
and Modularity of Objects



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## 3.5 Streams



実世界での課題: 変化 (change) のモデル化

1. 実世界での時間変化を計算オブジェクトの局所状態の時間変化 (time variation) をモデル化  
⇒ モデル化したオブジェクトの局所変数への代入 (assignment) で時間変化をとらえる
2. 代替案: stream を利用 ⇒ 一部の問題が軽減化  
瞬間での値変化ではなく, 値の全履歴 $x(t)$  で考える  
離散時間変化とみなすと,  $x(t)$  はsequence となる.  
Stream はsequence だが単なる list ではない.  
Delayed evaluation technique と組み合わせると,  
stream は大規模な(無限長の)sequence が表現可能

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### 3.5.1 Streams are delayed lists



2.2.3章: map, filter, accumulate, enumerate, ...

```
(define (sum-primes a b)
  (define (iter count accum)
    (cond ((> count b) accum)
          ((prime? count)
           (iter (+ count 1) (+ count accum)))
          (else (iter (+ count 1) accum))))
  (iter a 0))
```

簡潔な記述だが無駄な処理: 処理を順に適用, リストのコピーを繰り返す

```
(define (sum-primes a b)
  (accumulate
   +
   0
   (filter prime? (enumerate-interval a b))))
```

- (enumerate-interval 10000 1000000) は完全なリストを作る
- (car (cdr (filter prime? (enumerate-interval 10000 1000000))))

⇒ 10000以上の2番目の素数を返す, ただし, 全ての素数を求めてから.

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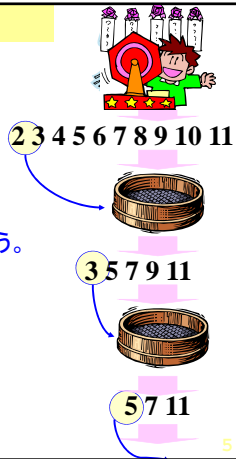
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### seq: 慣用インタフェース

- 処理間のインタフェース
- API (Application Program Interface)
- Parameterでの受け渡し
- データ構造をインタフェースに使う。
- sequence を活用
- 例: 素数を求めるための The Sieve of Eratosthenes (エラトステネスの篩)




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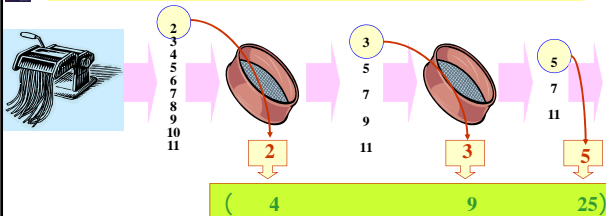
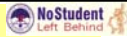
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### 共通性の視点: 素数の2乗を求める



共通点を見る4つの基本手続き

- 数え上げ(enumerate)
- フィルタ(filter)
- 写像(map)
- 集約(accumulate)




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## ストリーム: sequenceの簡潔表現+逐次実行

1. Formulate programs elegantly as sequence manipulation,
2. Attain the efficiency of incremental computation.

### 実装のアイデア:

- to arrange to construct a stream only partially and to pass the partial construction to the program that consumes the stream.
- If the consumer attempts to access a part of the stream that has not yet been constructed, the stream will automatically construct just enough more of itself to produce the required part.



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## データ構造 ストリームの構築子, 選択子



- cons-stream
- stream-car, stream-cdr
- the-empty-stream, stream-null?
- stream-ref
- stream-map, stream-for-each

```
(stream-car (cons-stream x y)) = x  
(stream-cdr (cons-stream x y)) = y
```

ただし, stream-car 評価時には y は評価されず遅延  
stream-cdr 評価時に y が評価される.

```
(define (stream-ref s n)  
  (if (= n 0)  
      (stream-car s)  
      (stream-ref (stream-cdr s) (- n 1))))
```

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## ストリームの基本演算



```
(define (stream-map proc s)  
  (if (stream-null? s)  
      the-empty-stream  
      (cons-stream  
        (proc (stream-car s))  
        (stream-map proc (stream-cdr s)))))  
  
(define (stream-for-each proc s)  
  (if (stream-null? s)  
      'done  
      (begin (proc (stream-car s))  
              (stream-for-each proc  
                                (stream-cdr s)))))  
  
(define (display-stream s)  
  (stream-for-each display-line s))  
  
(define (display-line x)  
  (newline)  
  (display x))
```

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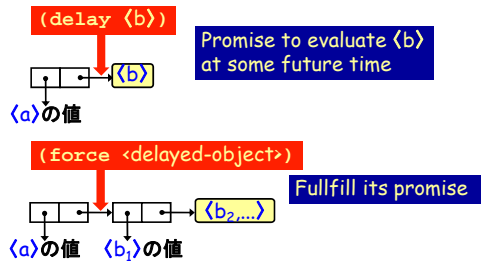
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## delay ストリーム 実装上のspecial form

- `(delay <exp>)`: delayed object を生成して返.
- `(force <obj>)`: delayed object を評価して結果を返す(fulfill its promise).

```
(cons-stream <a> <b>) = (cons <a>(delay <b>))
```



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## delay ストリーム 実装上のspecial form

```
(define (stream-car stream)  
  (car stream) )
```

```
(define (stream-cdr stream)  
  (force (cdr stream)) )
```

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## The stream implementaiton in action

### 2番目の素数の例の再考

```
(stream-car  
  (stream-cdr  
    (stream-filter prime?  
      (stream-enumerate-interval 10000 1000000) )))  
  
(define (stream-enumerate-interval low high)  
  (if (> low high)  
      the-empty-stream  
      (cons-stream  
        low  
        (stream-enumerate-interval (+ low 1) high) )))
```

### 実際の動きは

```
(stream-enumerate-interval 10000 1000000)  
⇒ (cons-stream 10000  
   (stream-enumerate-interval 10001 1000000) )  
⇒ (cons 10000  
   (delay (stream-enumerate-interval 10001 1000000)))
```

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**The stream implementation in action** 

```
(define (stream-filter pred stream)
  (cond ((stream-null? stream) the-empty-stream)
        ((pred (stream-car stream))
         (cons-stream (stream-car stream)
                       (stream-filter pred (stream-cdr stream))))
        (else (stream-filter pred (stream-cdr stream)))))
```

さあ、実行

10000~10006は素数でないので、stream-cdr をとってゆくと

```
(cons 10001
      (delay (stream-enumerate-interval 10002 1000000)))
```

....

```
(cons 10007
      (delay (stream-enumerate-interval 10008 1000000)))
```

10007 は素数なので、cons-stream を実行して

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**The stream implementation in action** 

```
(cons 10007
      (delay
        (stream-filter
          prime?
          (cons 10008
                (delay
                  (stream-enumerate-interval 10009
                    1000000 ))))))))
```

stream-cdr をとると

```
(cons 10009
      (delay
        (stream-filter
          prime?
          (cons 10010
                (delay
                  (stream-enumerate-interval 10011
                    1000000 ))))))))
```

stream-car をとると

```
10009
```

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**Implementing delay and force** 

(delay <exp>) を (lambda () <exp>) で表現すると

```
(define (force delayed-object)
  (delayed-object))
```

- 並列化では実行順序が不明 ⇒ 評価した結果を memoize

```
(define (memo-proc proc)
  (let ((already-run? false) (result false))
    (lambda ()
      (if (not already-run?)
          (begin (set! result (proc))
                 (set! already-run? true)
                 result)
          result ))))
```

- (delay <exp>) を (memo-proc (lambda () <exp>)) で表現

```
(define (force delayed-object)
  (delayed-object))
```

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### Exercise 3.50



Complete the following definition, which generalizes stream-map to allow procedures that take multiple arguments, analogous to map in section 2.2.3.

```
(define (stream-map proc s)
  (if (stream-null? s)
      the-empty-stream
      (cons-stream
        (proc (stream-car s))
        (stream-map proc (stream-cdr s)) )))

(define (stream-map proc . argstreams)
  (if (⟨??⟩ (car argstreams))
      the-empty-stream
      (⟨??⟩
        (apply proc (map ⟨??⟩ argstreams))
        (apply stream-map
          (cons proc (map ⟨??⟩ argstreams)) )))))
```

2引数を取る例を考えること

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### Exercise 3.51



In order to take a closer look at delayed evaluation, we will use the following procedure, which simply returns its argument after printing it:

```
(define (show x)
  (display-line x)
  x)
```

What does the interpreter print in response to evaluating each expression in the following sequence?

```
(define x (stream-map show
  (stream-enumerate-interval 0 10)))

(stream-ref x 5)

(stream-ref x 7)
```

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### Exercise 3.52



Consider the sequence of expressions

```
(define sum 0)
(define (accum x)
  (set! sum (+ x sum))
  sum)
(define seq (stream-map accum
  (stream-enumerate-interval 1 20)))
(define y (stream-filter even? seq))
(define z (stream-filter
  (lambda (x) (= (remainder x 5) 0)) seq))
(stream-ref y 7)
(display-stream z)
```

What is the value of sum after each of the above expressions is evaluated? What is the printed response to evaluating the stream-ref and display-stream expressions? Would these responses differ if we had implemented (delay <exp>) simply as (lambda () <exp>) without using the optimization provided by memo-proc? Explain.

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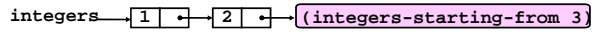
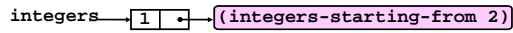
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### 3.5.2 Infinite Streams



- これまでのsequence は有限範囲.
- 無限ストリームを構築

```
(define (integers-starting-from n)
  (cons-stream n (integers-starting-from (+ n 1))))
(define integers (integers-starting-from 1))
```



```
(define (divisible? x y) (= (remainder x y) 0))
(define no-sevens
  (stream-filter (lambda (x) (not (divisible? x 7)))
                integers ))
```

```
(stream-ref no-sevens 100)
```

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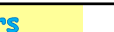
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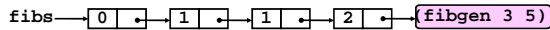
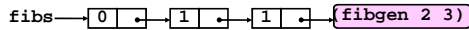
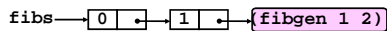
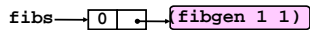
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### Infinite Stream of Fibonacci numbers



```
(define (fibgen a b)
  (cons-stream a (fibgen b (+ a b))))
(define fibs (fibgen 0 1))
```




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### Sieve of Eratosthenes



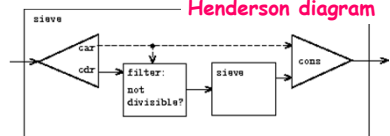
```
(define (sieve stream)
  (cons-stream
    (stream-car stream)
    (sieve (stream-filter
            (lambda (x)
              (not (divisible? x
                    (stream-car stream))))
            (stream-cdr stream))))))
```



```
(define primes
  (sieve (integers-starting-from 2)))
(stream-ref primes 50)
```

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Henderson diagram




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## Defining streams implicitly



```
(define ones (cons-stream 1 ones))
ones → 1 → ones
```

```
(define (add-streams s1 s2)
  (stream-map + s1 s2))
```

```
(define integers
  (cons-stream 1 (add-streams ones integers)))
```

integers → 1 → (add-streams ones integers)

integers → 1 → (stream-map + ones □)

integers → 1 → 2 → (stream-map + ones □)

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## Fibonacci number を stream で求める



```
(define (fib n)
  (cond ((= n 0) 0)
        ((= n 1) 1)
        (else (+ (fib (- n 1))
                  (fib (- n 2))))))
```

```
(define (fib-iter n)
  (define (iter a b count)
    (if (= count 0)
        b
        (iter (+ a b) a (- count 1))))
  (iter 1 0 n))
```

```
1 1 2 3 5 8 13 21 ... = (stream-cdr fibs)
0 1 1 2 3 5 8 13 ... = fibs
0 1 1 2 3 5 8 13 21 34 ... = fibs
```

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## Defining streams implicitly



```
(define fibs
  (cons-stream
    0
    (cons-stream
      1
      (add-streams (stream-cdr fibs) fibs))))
```

fibs → 0 → 1 → (add-streams (stream-cdr fibs) fibs)

fibs → 0 → 1 → (stream-map + □ □)

fibs → 0 → 1 → 1 → (stream-map + □ □)

fibs → 0 → 1 → 1 → 2 → (stream-map + □ □)

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## Scale streams



```
(define (scale-stream stream factor)
  (stream-map
    (lambda (x) (* x factor))
    stream ))
```

2 のべき乗の列(1, 2, 4, 8, 16, 32, ....)を生成

```
(define double
  (cons-stream 1 (scale-stream double 2)))
```

double → 1 → (scale-stream double 2)

double → 1 → 2 → (scale-stream 2)

double → 1 → 2 → 4 → (scale-stream 2)

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## Infinite Stream of Primes



```
(define primes
  (cons-stream
    2
    (stream-filter
      prime?
      (integers-starting-from 3))))
```

prime? の定義は難しい

```
(define (prime? n)
  (define (iter ps)
    (cond ((> (square (stream-car ps)) n)
           true)
          ((divisible? n (stream-car ps))
           false)
          (else (iter (stream-cdr ps)))))
  (iter primes))
```

primes と prime? は再帰的定義

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## Exercise 3.53



Without running the program, describe the elements of the stream defined by

```
(define s
  (cons-stream 1 (add-streams s s)))
```

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### Exercise 3.54



Define a procedure `mul-streams`, analogous to `add-streams`, that produces the elementwise product of its two input streams. Use this together with the stream of integers to complete the following definition of the stream whose  $n$ th element (counting from 0) is  $n + 1$  factorial:

$n! = 1 * 2 * 3 * \dots$  で定義する

```
(define factorials
  (cons-stream 1
    (mul-streams <??> <??> ) )
```

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### 宿題: 5月23日午前8時 締切



1. Ex3.50(実行例を考える), 3.51, 3.52, 3.54.
  2. プログラムの説明, 実行例をつけて, 設問に応えること.
  3. レポート(PDF)とプログラムファイルを送付  
**PROG-6@zeus.kuis.kyoto-u.ac.jp**  
**file 名は 学籍番号-名前-6.pdf**
- 友達に教えてもらったら, その人の名前を明記すること. Webは出展を明記. (otherwise 『同じ』回答は減点)



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